

## Content-Guidance in Formal Problem Solving Processes

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### Some Background (1)

- ▶ Vienna Circle: (degenerated to) a priori methodology
- ▶ Historicism (Kuhn, ...): relativism
- ▶ today: scientific problem-solving is **content-guided** "what we have learned, including what we have learned about how to learn" (Shapere)

all possibilities / historicism

Aim of this paper:

- ▶ formal approach to problem solving
- ▶ in which there is large room for content-guidance

(⇒ example)

### Some Background (3)

some elements of the plot:

- ▶ aim: *explication* for problem solving processes (psps)
- ▶ *backbone*: solve  $\{?A, \sim A\}$  by **CL**-deriving  $A$  or  $\sim A$  from  $\Gamma$
- ▶ empirical means: observation and experiment guided by psp
- ▶ new available information (not originally seen as relevant) guided by psp
- ▶ adaptive logics:
  - corrective*: handling inconsistency (and similar), ambiguity, vagueness, ...
  - ampliative*: inductive generalization, abduction, ...
- ▶ control defeasible inferences by conditions and marking
- ▶ model-based reasoning (mainly future research)

### Elements of a psp backbone

lines of psp

- ▶ problem lines: example  $\{?{p \vee q, \sim p \vee q}, ?{r, \sim r}\}$   
problem = non-empty set of questions
- ▶ declarative lines
  - conditional:  $[B_1, \dots, B_n] A$
  - unconditional:  $[\emptyset] A$ , viz.  $A$

### Outline

Some Background

Elements of a psp backbone

- procedure
- rules
- instructions
- marking

An Example of the Backbone

Adaptive logics

Some Extra-Logical Extensions

Comments

### Some Background (2)

direct sources of inspiration:

- ▶ philosophy of science/epistemology:
  - ▶ problem solving: Kuhn, Laudan, *Nickles*, ...
  - ▶ 'contextual': specific certainties, relevant data, methodological instructions, ... for each problem
- ▶ logic:
  - ▶ procedures
  - ▶ prospective dynamics
  - ▶ adaptive logics
  - ▶ erotetic logic (varying on Wiśniewski)

problem determined by (changing) *constraints*

- ▶ conditions on solution
- ▶ methodological instructions / heuristics / examples
- ▶ certainties (conceptual system, ...)

### Some Background (4)

from the general plot (consolutions?)

- ▶ in the end all knowledge defeasible
- ▶ going through different 'contexts'
  - ⇒  $\left\{ \begin{array}{l} \text{conceptual change} \\ \text{generation of new concepts} \end{array} \right.$
- ▶ many ampliative mechanisms (abduction overestimated)
- ▶ independent of ontological debates

psp

- ▶ stage of a psp: sequence of lines
- ▶ psp: chain of stages
- ▶ next stage: add new line + apply marking definitions
- ▶ governed by *procedure*

*procedure* (set of instructions)

- ▶ *rule of inference*: to derive  $A$  from  $B_1, \dots, B_n$
- ▶ *instruction*:  
rule of inference + permissions/obligations  
permissions/obligations depend on the present stage (lines + marks)
- ▶ below: instructions (including the rules)

varying on Smullyan

a	a <sub>1</sub>	a <sub>2</sub>	b	b <sub>1</sub>	b <sub>2</sub>
$A \wedge B$	$A$	$B$	$\sim(A \wedge B)$	$*A$	$*B$
$A \equiv B$	$A \supset B$	$B \supset A$	$\sim(A \equiv B)$	$\sim(A \supset B)$	$\sim(B \supset A)$
$\sim(A \vee B)$	$*A$	$*B$	$A \vee B$	$A$	$B$
$\sim(A \supset B)$	$A$	$*B$	$A \supset B$	$*A$	$B$
$\sim\sim A$	$A$	$A$			

complement of A:  $*A = B$  if  $A = \sim B$ ; otherwise  $*A = \sim A$

a	a <sub>1</sub>	a <sub>2</sub>	b	b <sub>1</sub>	b <sub>2</sub>
$A \wedge B$	$A$	$B$	$\sim(A \wedge B)$	$*A$	$*B$
$A \equiv B$	$A \supset B$	$B \supset A$	$\sim(A \equiv B)$	$\sim(A \supset B)$	$\sim(B \supset A)$
$\sim(A \vee B)$	$*A$	$*B$	$A \vee B$	$A$	$B$
$\sim(A \supset B)$	$A$	$*B$	$A \supset B$	$*A$	$B$
$\sim\sim A$	$A$	$A$			

positive part relation

1.  $pp(A, A)$ .
2.  $pp(A, a)$  if  $pp(A, a_1)$  or  $pp(A, a_2)$ .
3.  $pp(A, b)$  if  $pp(A, b_1)$  or  $pp(A, b_2)$ .
4. If  $pp(A, B)$  and  $pp(B, C)$ , then  $pp(A, C)$ .

Formula analysing rules:

$$\frac{[\Delta] \alpha}{[\Delta] \alpha_1 \quad [\Delta] \alpha_2} \quad \frac{[\Delta] \beta}{[\Delta \cup \{*\beta_2\}] \beta_1 \quad [\Delta \cup \{*\beta_1\}] \beta_2}$$

FAR If  $C$  is an unmarked target,  $[\Delta] A$  is the formula of an unmarked line  $i$ ,  $[\Delta] A / [\Delta \cup \Delta'] B$  is a formula analysing rule, and  $pp(C, B)$ , then one may add:

$$k \quad [\Delta \cup \Delta'] B \quad i; R$$

in which  $R$  is the name of the formula analysing rule.

Eliminate some problems without answering them:

EM0 If  $[\Delta \cup \{*\alpha\}] A$  is the formula of a line  $i$  that is neither R-marked nor I-marked, then one may add:

$$k \quad [\Delta] A \quad i; EM0$$

EM If  $A$  is an unmarked target,  $[\Delta \cup \{B\}] A$  and  $[\Delta' \cup \{\sim B\}] A$  are the respective formulas of the unmarked or only D-marked lines  $i$  and  $j$ , and  $\Delta \subseteq \Delta'$  or  $\Delta' \subseteq \Delta$ , then one may add:

$$k \quad [\Delta \cup \Delta'] A \quad i, j; EM$$

a	a <sub>1</sub>	a <sub>2</sub>	b	b <sub>1</sub>	b <sub>2</sub>
$A \wedge B$	$A$	$B$	$\sim(A \wedge B)$	$*A$	$*B$
$A \equiv B$	$A \supset B$	$B \supset A$	$\sim(A \equiv B)$	$\sim(A \supset B)$	$\sim(B \supset A)$
$\sim(A \vee B)$	$*A$	$*B$	$A \vee B$	$A$	$B$
$\sim(A \supset B)$	$A$	$*B$	$A \supset B$	$*A$	$B$
$\sim\sim A$	$A$	$A$			

formula analysing rules:

$$\frac{[\Delta] \alpha}{[\Delta] \alpha_1 \quad [\Delta] \alpha_2} \quad \frac{[\Delta] b}{[\Delta \cup \{*\beta_2\}] b_1 \quad [\Delta \cup \{*\beta_1\}] b_2}$$

condition analysing rules:

$$\frac{[\Delta \cup \{a\}] A}{[\Delta \cup \{a_1, a_2\}] A} \quad \frac{[\Delta \cup \{b\}] A}{[\Delta \cup \{b_1\}] A \quad [\Delta \cup \{b_2\}] A}$$

The instructions

Main Start a psp with the line:

$$1 \quad \{?\{M, \sim M\}\} \quad \text{Main}$$

Target If  $P$  is the problem of an unmarked problem line, and  $A$  is a direct answer of a member of  $P$ , then one may add:

$$k \quad [A] A \quad \text{Target}$$

Prem If  $A$  is an unmarked target,  $B \in \Gamma$ , and  $pp(A, B)$ , then one may add:

$$k \quad B \quad \text{Prem}$$

Condition analysing rules:

$$\frac{[\Delta \cup \{a\}] A}{[\Delta \cup \{a_1, a_2\}] A} \quad \frac{[\Delta \cup \{b\}] A}{[\Delta \cup \{b_1\}] A \quad [\Delta \cup \{b_2\}] A}$$

CAR If  $A$  is an unmarked target,  $[\Delta \cup \{B\}] A$  is the formula of an unmarked line  $i$ , and  $[\Delta \cup \{B\}] A / [\Delta \cup \Delta'] A$  is a condition analysing rule, then one may add:

$$k \quad [\Delta \cup \Delta'] A \quad i; R$$

in which  $R$  is the name of the condition analysing rule.

eliminate solved questions from a problem and summarize remaining problems (and paths):

Trans If  $A$  is an unmarked target, and  $[\Delta \cup \{B\}] A$  and  $[\Delta'] B$  are the respective formulas of the at most S-marked lines  $i$  and  $j$ , then one may add:

$$k \quad [\Delta \cup \Delta'] A \quad i, j; \text{Trans}$$

handle derived problems:

DP If  $A$  is an unmarked target from problem line  $i$  and  $[B_1, \dots, B_n]$   $A$  is the formula of an unmarked line  $j$ , then one may add:  
 $k \quad \{?\{B_1, \sim B_1\}, \dots, ?\{B_n, \sim B_n\}\} \quad i, j; DP$

Remark:

no instruction for applying EFQ  
in view of the intended applications  
(deriving predictions, explanations, etc.)

### Marking definitions

redundant lines are R-marked:

#### Definition

An at most S-marked declarative line  $i$  that has  $[\Delta] A$  as its formula is R-marked at a stage iff, at that stage,  $[\Theta] A$  is the formula of a line for some  $\Theta \subset \Delta$ .

#### Definition

An unmarked problem line  $i$  is R-marked at a stage iff, at that stage, a direct answer  $A$  of a question of line  $i$  is the formula of a line.

- target from a problem line
- resolution line
- direct target from
- target sequence
- grounded target sequence

inoperative lines are I-marked (not useful for extant problem):

#### Definition

An at most S-marked target line that has  $[A] A$  as its formula is I-marked at a stage iff every problem line from which  $A$  is a target is marked at that stage.

#### Definition

An at most S-marked resolution line of which  $[\Delta^1] A^1$  is the formula and  $\Delta^1 \neq \emptyset$  is I-marked at a stage iff, at that stage, for every grounded target sequence  $\langle [\Delta^n] A^n, \dots, [\Delta^1] A^1 \rangle$ ,

- (i) some target  $[A^i] A^i$  ( $1 \leq i \leq n$ ) is marked, or
- (ii)  $\{A^n, \dots, A^1\} \cap \Delta^1 \neq \emptyset$ , or
- (iii)  $\Delta^1 \cup \dots \cup \Delta^n \cup \Gamma_s^o$  is flatly inconsistent.

#### Definition

An unmarked problem line is I-marked iff no unmarked resolution line generates it.

Dead end lines are D-marked (no further action from such line)

- $A$  is a *dead end* ( $A$  is literal and not positive part of premise)
- CAR-descendant of  $[\Delta] B$

#### Definition

An at most S-marked resolution line with formula  $[\Delta] A$  is D-marked at a stage iff some  $B \in \Delta$  is a dead end or, at that stage, all CAR-descendants of  $[\Delta] A$  occur in the psp and are D-marked.

#### Definition

An at most S-marked target line with formula  $[A] A$  is D-marked at a stage iff  $A$  is a dead end or no further action can be taken in view of target  $A$ .

for all consistent  $\Gamma$ :

if  $\Gamma \vdash A$ , then the procedure applied to  $\Gamma$  and  $\{?\{A, \sim A\}\}$  results in the answer  $A$ ,

and

if  $\Gamma \not\vdash A$ , then the procedure applied to  $\Gamma$  and  $\{?\{A, \sim A\}\}$  stops without the main problem being answered or results in the answer  $\sim A$ .

Speed up the procedure by S-marks

- $\Gamma_s^o$  union of the  $\Gamma$  and of the set of the conditionless formulas that occur at stage  $s$  of the psp

#### Definition

A R-unmarked resolution line in which  $[\Delta^1] A^1$  is derived is S-marked iff

- (i)  $\Delta^1 \cap \Gamma_s^o \neq \emptyset$ , or
- (ii) for some target sequence  $\langle [\Delta^n] A^n, \dots, [\Delta^1] A^1 \rangle$ ,  $\{A^n\} \cup \Delta^1$  is flatly inconsistent whereas  $\Delta^1$  is not flatly inconsistent, or
- (iii)  $\Delta^1 \subset \Delta^n \cup \dots \cup \Delta^2$  for some target sequence  $\langle [\Delta^n] A^n, \dots, [\Delta^1] A^1 \rangle$ .

instruction: operate on S-marked lines before doing anything else

## An Example of the Backbone

main problem:  $\{p \vee q, \sim(p \vee q)\}$

premise set:  $\{\sim s, \sim u \vee r, (r \wedge t) \vee s, (q \vee u) \supset (\sim t \vee q), t \supset u\}$

logic:  $\mathbf{CL}^-$  + an erotetic logic (fixed by the procedure)

procedure: prospective dynamics + problems

$\{\sim s, \sim u \vee r, (r \wedge t) \vee s, (q \vee u) \supset (\sim t \vee q), t \supset u\}$

1	$\{? \{p \vee q, \sim(p \vee q)\}\}$	Main	
...			
6	$[q] p \vee q$	4; CVE	
...			
14	$[u, t] q$	13; VE	S <sup>23</sup> R <sup>24</sup>
15	$\{? \{u, \sim u\}, ? \{t, \sim t\}\}$	8, 14; DP	R <sup>23</sup>
16	$[t] t$	Target	R <sup>23</sup>
17	$(r \wedge t) \vee s$	Prem	
18	$[\sim s] r \wedge t$	17; VE	
19	$[\sim s] t$	18; $\wedge$ E	S <sup>22</sup> R <sup>23</sup>
20	$\{? \{s, \sim s\}\}$	16, 19; DP	R <sup>22</sup>
21	$[\sim s] \sim s$	Target	R <sup>22</sup>
22	$\sim s$	Prem	
23	$t$	19, 22; Trans	
24	$[u] q$	14, 23; Trans	
25	$\{? \{u, \sim u\}\}$	8, 24; DP	

1, 6, 23–25

$\{\sim s, \sim u \vee r, (r \wedge t) \vee s, (q \vee u) \supset (\sim t \vee q), t \supset u\}$

1	$\{? \{p \vee q, \sim(p \vee q)\}\}$	Main	
2	$[\sim(p \vee q)] \sim(p \vee q)$	Target	D <sup>3</sup>
3	$[\sim p, \sim q] \sim(p \vee q)$	2; C $\sim$ VE	D <sup>3</sup>
4	$[p \vee q] p \vee q$	Target	
5	$[p] p \vee q$	4; CVE	D <sup>5</sup>
6	$[q] p \vee q$	4; CVE	
7	$\{? \{q, \sim q\}\}$	4, 6; DP	
8	$[q] q$	Target	
9	$(q \vee u) \supset (\sim t \vee q)$	Prem	
10	$[q \vee u] \sim t \vee q$	9; $\supset$ E	
11	$[q] \sim t \vee q$	10; CVE	
12	$[q, t] q$	11; VE	I <sup>12</sup>
13	$[u] \sim t \vee q$	10; CVE	
14	$[u, t] q$	13; VE	
15	$\{? \{u, \sim u\}, ? \{t, \sim t\}\}$	8, 14; DP	
16	$[t] t$	Target	

1, 6, 14–16

$\{\sim s, \sim u \vee r, (r \wedge t) \vee s, (q \vee u) \supset (\sim t \vee q), t \supset u\}$

1	$\{? \{p \vee q, \sim(p \vee q)\}\}$	Main	R <sup>31</sup>
...			
6	$[q] p \vee q$	4; CVE	S <sup>30</sup> R <sup>31</sup>
...			
23	$t$	19, 22; Trans	
24	$[u] q$	14, 23; Trans	S <sup>29</sup> R <sup>30</sup>
25	$\{? \{u, \sim u\}\}$	8, 24; DP	R <sup>29</sup>
26	$[u] u$	Target	R <sup>29</sup>
27	$t \supset u$	Prem	
28	$[t] u$	27; $\supset$ E	S <sup>28</sup> R <sup>29</sup>
29	$u$	23, 28; Trans	
30	$q$	24, 29; Trans	
31	$p \vee q$	6, 30; Trans	

problem solved

## Adaptive logics (only Standard Format)

characterization

- ▶ lower limit logic **LLL**  
monotonic, compact, ... logic
- ▶ set of abnormalities  $\Omega$   
characterized by a (possibly restricted) logical form
- ▶ strategy  
Reliability, Minimal Abnormality, ...

upper limit logic:

syntax:  $\mathbf{ULL} = \mathbf{LLL} + \text{axiom/rule}$   
semantics: the **LLL**-models that verify no abnormality

general idea behind adaptive logics:

$Cn_{\mathbf{AL}}(\Gamma) : Cn_{\mathbf{LLL}}(\Gamma) + \text{what follows if as many members of } \Omega \text{ are false as the premises permit}$

Example: logic of inductive generalization: **IL<sup>m</sup>**

- ▶ lower limit logic: **CL**
- ▶ set of abnormalities:  $\Omega = \{\exists A \wedge \exists \sim A \mid A \in \mathcal{F}^0\}$
- ▶ strategy: Minimal Abnormality

upper limit logic:

$\mathbf{UCL} = \mathbf{CL} + \exists \alpha A(\alpha) \supset \forall \alpha A(\alpha)$

semantically: the uniform **CL**-models  $(v(\pi^r) \in \{\emptyset, D^{(r)}\})$

ampliative adaptive logic (if **CL** is the standard)

Example: the inconsistency-adaptive **CLuN<sup>r</sup>**

- ▶ lower limit logic: **CLuN**
- ▶ set of abnormalities:  $\Omega = \{\exists(A \wedge \sim A) \mid A \in \mathcal{F}\}$
- ▶ strategy: Reliability

upper limit logic:

$\mathbf{CL} = \mathbf{CLuN} + (A \wedge \sim A) \supset B$

semantically: the **CLuN**-models that verify no inconsistency

corrective adaptive logic (if **CL** is the standard)

standard format provides

- ▶ proofs
- ▶ semantics
- ▶ most of metatheory (including soundness and completeness)
- ▶ prospective dynamics (published for Reliability)

further examples (relevant for philosophy of science)

- ▶ many other inconsistency-handling (+ other logical symbols)
- ▶ ambiguity-adaptive
- ▶ vagueness-adaptive
- ▶ corrective deontic logics
- ▶ paraconsistent compatibility
- ▶ ...
- ▶ plausibility-adaptive
- ▶ compatibility
- ▶ diagnosis
- ▶ abduction
- ▶ analogies, metaphors
- ▶ erotetic evocation/implication (problem solving)
- ▶ ...

## Some Extra-Logical Extensions (1)

answerable questions

$\mathbb{A}$  is a set of couples  $(\Delta : Q)$  in which  $\Delta$  is a set of statements and  $Q$  is a question

idea: if the members of  $\Delta$  are true,  $Q$  can be answered by observational/experimental means (not Hintikka's oracle)

New If  $A$  is an unmarked target,  $pp(A, B)$  for some direct answer  $B$  of  $Q$ ,  $(\Delta : Q)$  and all members of  $\Delta$  occur in the fsp, then one may add, for some direct answer  $C$  of  $Q$ :

$k \quad C \quad i$ ; New

psp guides (which observations/experiments should be carried out)

## Some Extra-Logical Extensions (2)

bringing in available information (formerly judged irrelevant)

one tries to solve problem from theory  $T$  and set of data  
later a theory  $T'$  turns out to be relevant  
(because a target is a positive part of an axiom of  $T'$ )

psp guides (which further theories are relevant?)

## Some Extra-Logical Extensions (3)

plausible conjectures

where  $A$  is an abnormality, introduce  $\diamond^i A$  or  $\diamond^i \neg A$

- ▶ basis: worldview, personal constraint, study of situation, blind guess
- ▶ thus reducing a disjunction of abnormalities
- ▶ = defeasibly obtaining more consequences (plausibility-adaptive logic)

psp guides (which disjunctions of abnormalities may be reduced?)

...

## Comments (1)

framework that contains **open slots**

these make content guidance possible

but the framework is **formal**

prospective dynamics pushes the 'logical' part of the heuristics into the proof

part of remaining heuristics is fixed by procedure

still remaining heuristics

## Comments (2)

content-guided

- ▶ 'language' of a scientific discipline (not typical)
- ▶ adaptive logics validate *applications* of rules that transcend the lower limit logic
- ▶ multiplicity of adaptive logics for every purpose (to be justified)
- ▶ multiplicity of erotetic logics (to be justified)
- ▶ multiplicity of procedures for prospective dynamics (to be justified)
- ▶ take background theories serious + several forms of defeasibility
- ▶ ./.

## Comments (3)

content-guided

- ▶ ./.
  - ▶ 'guesses': world-view, personal constraints, ..., blind (which guesses useful: determined by disjunctions of abnormalities) (extra logical origin; logic guides handling of the guesses)
  - ▶ local selection of adaptive logics (abd./ind.; inconsistency; replace lower limit logic; plausibilities; ...)
  - ▶ heuristics of psp
    - ▶ road followed to derive conclusion
    - ▶ observation / experiment / theoretical derivation
    - ▶ [use of models]
- to be decided in view of what was learned about world/learning in specific domain/context

## Comments (4)

conclusion

framework that contains **open slots**

these make content guidance possible

but the framework is **formal**

status of the approach itself: provisional hypothesis